

Collision Attacks on Up to 5 Rounds of SHA-3 Using Generalized Internal Differentials

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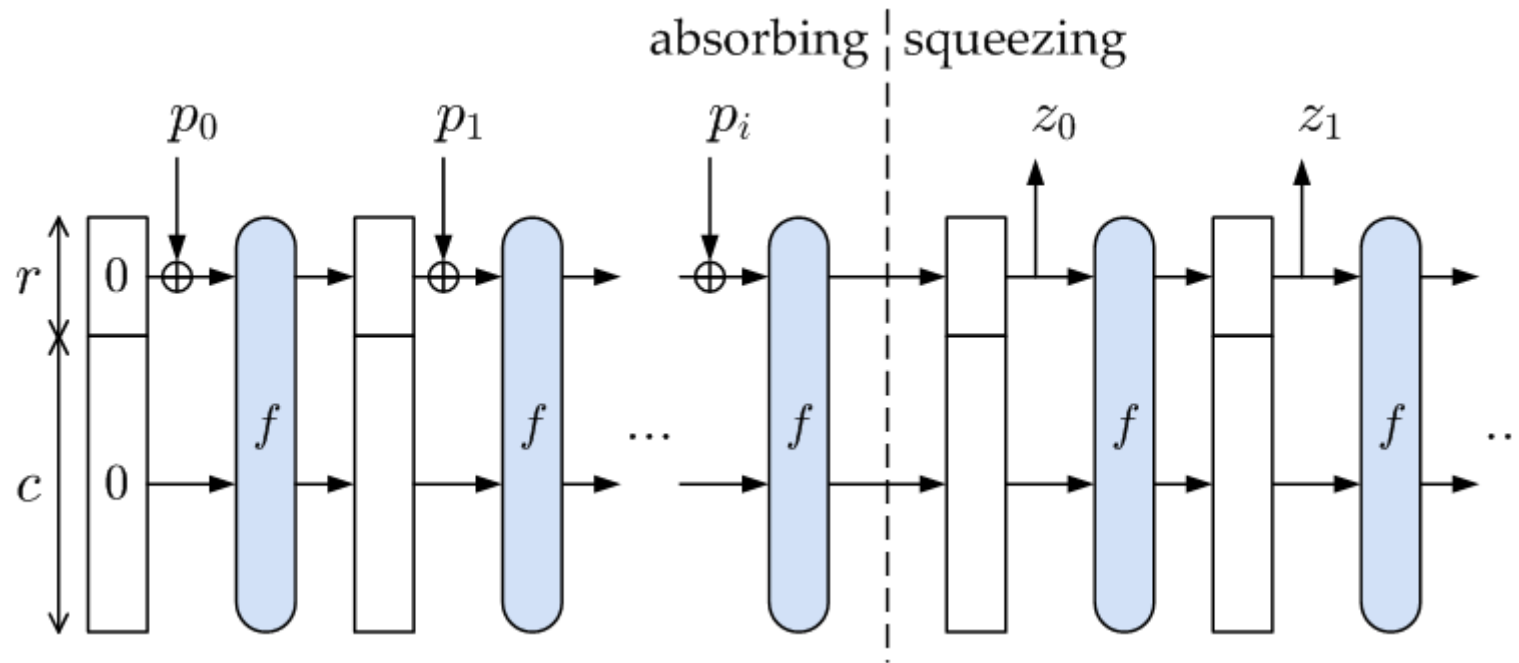
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Keccak

(Bertoni, Daemen, Peeters and Van Assche)

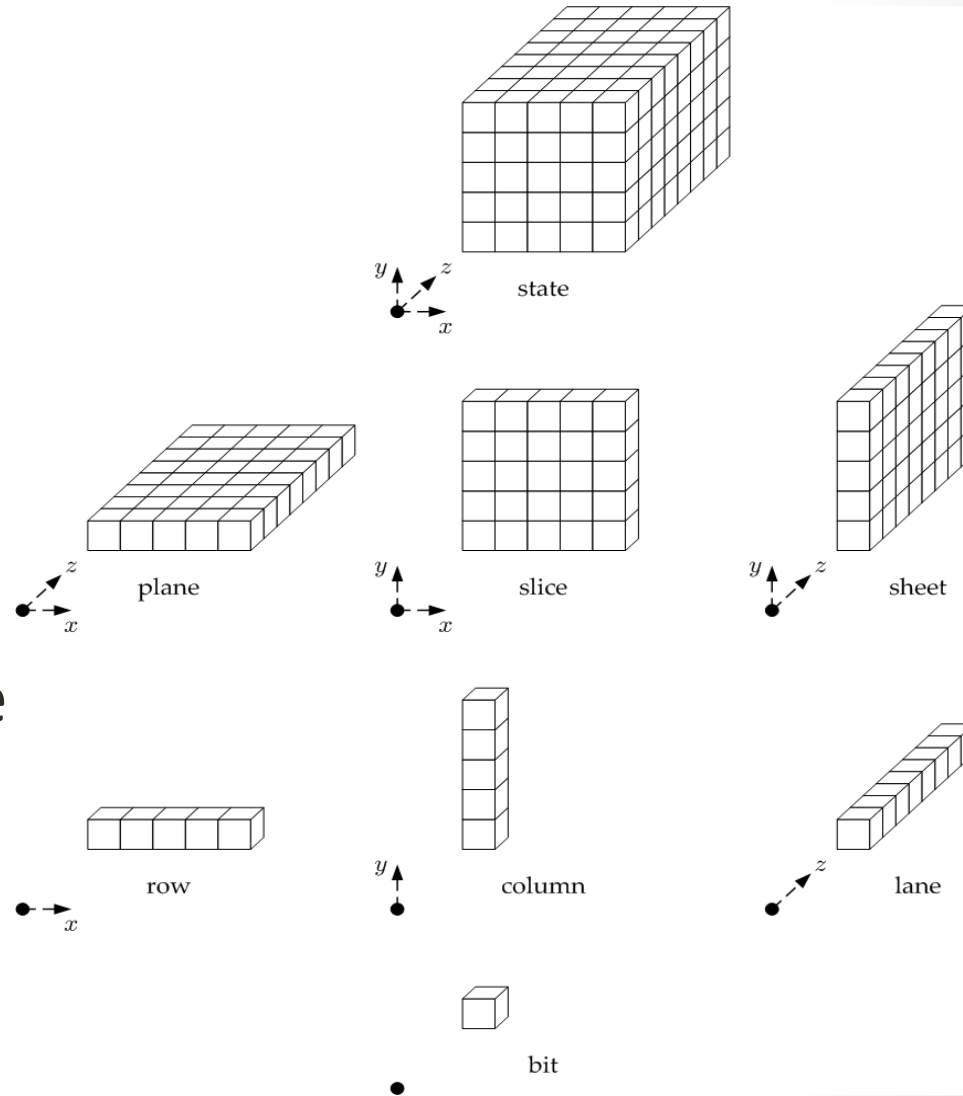
- Uses the **sponge construction**



Keccak

The Inner State

- Can be viewed as a **5x5x64**-bit cube
- Or as a **5x5** matrix, where each cell is a **64**-bit lane in the direction of the **z** axis



Keccak

The function f

- f is a **24-round** permutation on the 1600-bit state
- Each round consists of **5** mappings $R = \iota \circ \chi \circ \pi \circ \rho \circ \Theta$
- We denote $L = \pi \circ \rho \circ \Theta$ and refer to L as a “**half-round**”, where $\iota \circ \chi$ make up the other half

Keccak

The function f

- χ is the only **non-linear** mapping of Keccak
 - It has an algebraic degree of **2**
- ι adds a **low Hamming-weight** round constant to the state
- The state is initialized to **zero** before the XOR with the first message block

Keccak

Collision Attacks on Round-Reduced Keccak

- “Practical analysis of reduced-round Keccak” by Naya-Plasencia, Röck and Meier (Indocrypt 2011)
 - Collisions in **2** rounds of **Keccak-224** and **Keccak-256**
- “New attacks on Keccak-224 and Keccak-256” by Dinur, Dunkelman and Shamir (FSE 2012)
 - Collisions in **4** rounds of **Keccak-224** and **Keccak-256**
- **No** published collision attack on **Keccak-384** and **Keccak-512**

Keccak

Our New Results

- **Keccak-512**: A 3-round **practical** collision attack
- **Keccak-384**: A 3-round **practical** collision attack
- A 4-round collision attack (faster than the birthday bound by 2^{45})
- **Keccak-256**: A 5-round collision attack (faster than the birthday bound by 2^{13})

	Keccak-224	Keccak-256	Keccak-384	Keccak-512
Previous	4 (practical)	4 (practical)	-	-
New	-	5 (2^{115})	3 (practical) 4 (2^{147})	3 (practical)

Keccak

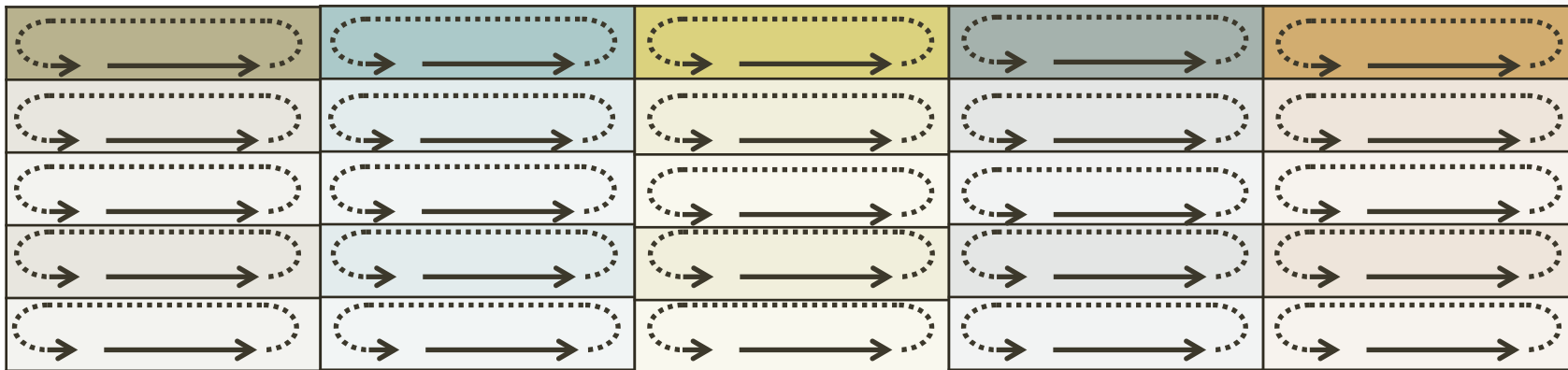
The Translation-Invariance Property

- Defined in the Keccak submission document
- 4 out of the 5 internal mappings (all but ι) are translation invariant in the direction of the z axis (of length 64)

Keccak

The Translation-Invariance Property

- If one state is the rotation of the other with respect to the z-axis, then applying to them any of the Θ, ρ, π, χ operations, **maintains this property**



Symmetric States

- A state which is **rotation-invariant** in the direction of the **z** axis by some **rotation index i** is called a **symmetric state**
- **i** can attain non-trivial values that divide the lane size **64** ($i \in \{1, 2, 4, 8, 16, 32\}$)

Consecutive Slice Sets

An example

- For $i=16$ we split the state into **4 consecutive slice sets (CSS)**

a₁				b₁				c₁				d₁				e₁			
f ₁				g ₁				h ₁				i ₁				j ₁			
k ₁				l ₁				m ₁				n ₁				o ₁			
p ₁				q ₁				r ₁				s ₁				t ₁			
u ₁				v ₁				w ₁				x ₁				y ₁			

	a₂				b₂				c₂				d₂				e₂			
	f ₂				g ₂				h ₂				i ₂				j ₂			
	k ₂				l ₂				m ₂				n ₂				o ₂			
	p ₂				q ₂				r ₂				s ₂				t ₂			
	u ₂				v ₂				w ₂				x ₂				y ₂			

Symmetric States

An Example

- In symmetric states all CSS's are equal
- In a symmetric state with $i=16$, each 64-bit lane is composed of a 4-repetition of a 16-bit value

a_1	a_1	a_1	a_1	b_1	b_1	b_1	b_1	c_1	c_1	c_1	c_1	d_1	d_1	d_1	d_1	e_1	e_1	e_1	e_1
f_1	f_1	f_1	f_1	g_1	g_1	g_1	g_1	h_1	h_1	h_1	h_1	i_1	i_1	i_1	i_1	j_1	j_1	j_1	j_1
k_1	k_1	k_1	k_1	l_1	l_1	l_1	l_1	m_1	m_1	m_1	m_1	n_1	n_1	n_1	n_1	o_1	o_1	o_1	o_1
p_1	p_1	p_1	p_1	q_1	q_1	q_1	q_1	r_1	r_1	r_1	r_1	s_1	s_1	s_1	s_1	t_1	t_1	t_1	t_1
u_1	u_1	u_1	u_1	v_1	v_1	v_1	v_1	w_1	w_1	w_1	w_1	x_1	x_1	x_1	x_1	y_1	y_1	y_1	y_1

Symmetric states remain symmetric after applying the Θ, ρ, π, χ operations

a_1	a_1	a_1	a_1	b_1	b_1	b_1	b_1	c_1	c_1	c_1	c_1	d_1	d_1	d_1	d_1	e_1	e_1	e_1	e_1
f_1	f_1	f_1	f_1	g_1	g_1	g_1	g_1	h_1	h_1	h_1	h_1	i_1	i_1	i_1	i_1	j_1	j_1	j_1	j_1
k_1	k_1	k_1	k_1	l_1	l_1	l_1	l_1	m_1	m_1	m_1	m_1	n_1	n_1	n_1	n_1	o_1	o_1	o_1	o_1
p_1	p_1	p_1	p_1	q_1	q_1	q_1	q_1	r_1	r_1	r_1	r_1	s_1	s_1	s_1	s_1	t_1	t_1	t_1	t_1
u_1	u_1	u_1	u_1	v_1	v_1	v_1	v_1	w_1	w_1	w_1	w_1	x_1	x_1	x_1	x_1	y_1	y_1	y_1	y_1

↓ Θ, ρ, π, χ

a_2	a_2	a_2	a_2	b_2	b_2	b_2	b_2	c_2	c_2	c_2	c_2	d_2	d_2	d_2	d_2	e_2	e_2	e_2	e_2
f_2	f_2	f_2	f_2	g_2	g_2	g_2	g_2	h_2	h_2	h_2	h_2	i_2	i_2	i_2	i_2	j_2	j_2	j_2	j_2
k_2	k_2	k_2	k_2	l_2	l_2	l_2	l_2	m_2	m_2	m_2	m_2	n_2	n_2	n_2	n_2	o_2	o_2	o_2	o_2
p_2	p_2	p_2	p_2	q_2	q_2	q_2	q_2	r_2	r_2	r_2	r_2	s_2	s_2	s_2	s_2	t_2	t_2	t_2	t_2
u_2	u_2	u_2	u_2	v_2	v_2	v_2	v_2	w_2	w_2	w_2	w_2	x_2	x_2	x_2	x_2	y_2	y_2	y_2	y_2

The Fifth Mapping

- **l** destroys the perfect symmetry of the state by adding a non-symmetric round constant

An Overview of the Basic Attack

- Pick a single-block message such that the initial state is **symmetric**
- The state **remains symmetric** after the first **4** mappings
- The symmetry is **slightly perturbed** by the **ι** mapping since the constants added are of **low Hamming-weight** (between **1** and **5**)
- The **diffusion** is sufficiently **slow** such that the state remains “close” to symmetric for the first few rounds

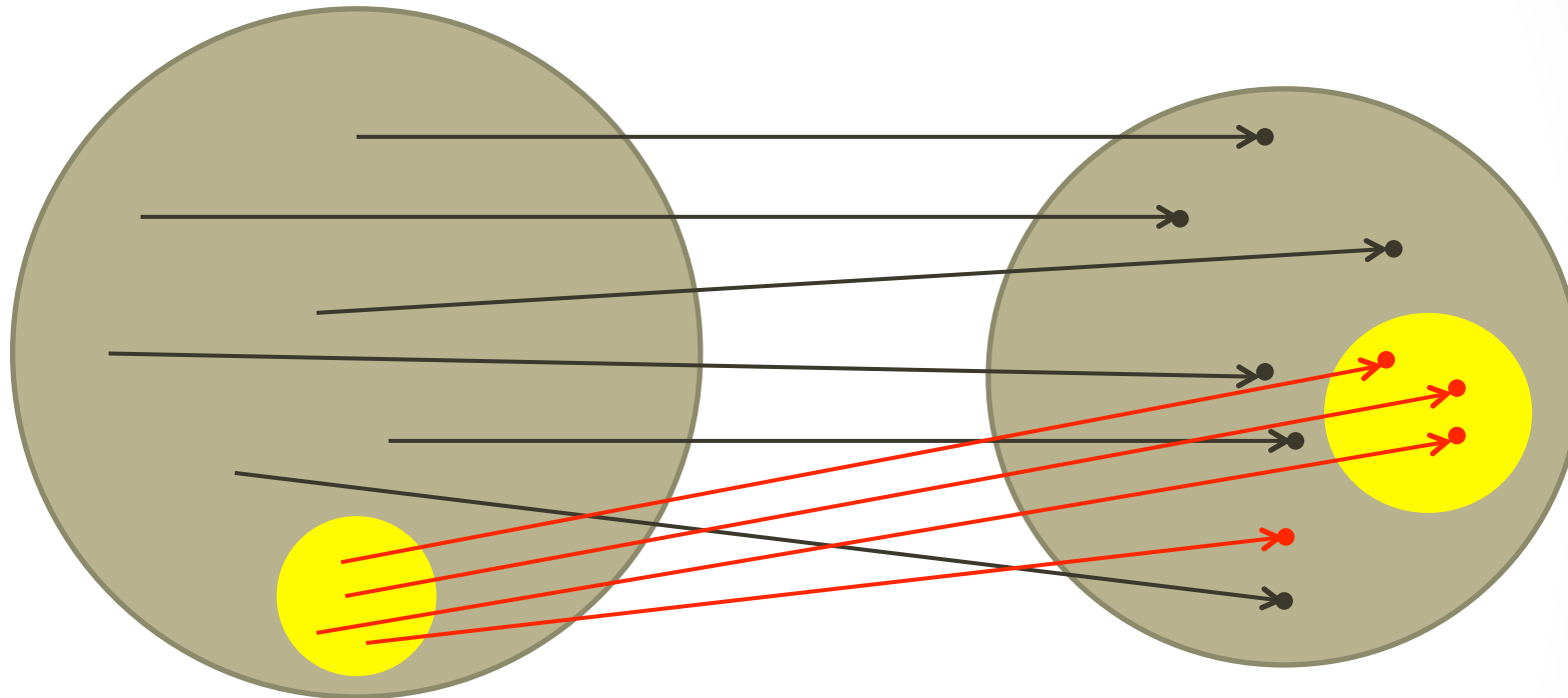
An Overview of the Basic Attack

The Squeeze Attack

- The **effective output size** for symmetric messages is **reduced**
- We use a natural attack (called the **squeeze attack**) that exploits this property
- We force a larger than expected number of inputs to squeeze into a **small subset** of possible outputs in which collisions are **more likely**

An Overview of the Basic Attack

The Squeeze Attack



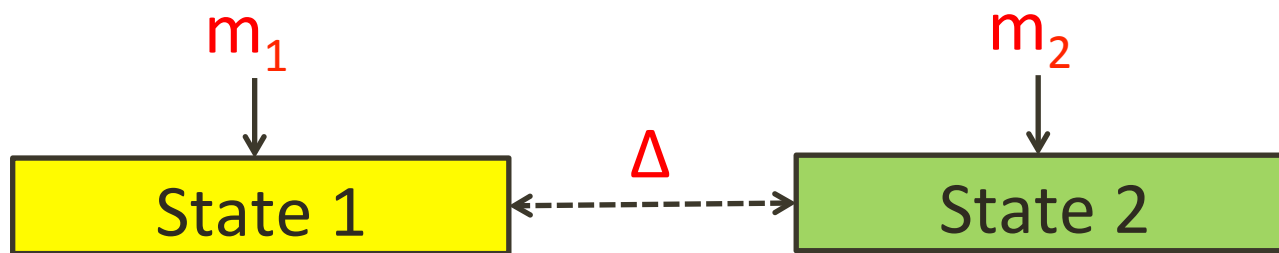
- A member of the input set is mapped with probability p to the output set of size D
- The time complexity of the attack is $1/p \cdot \nu D$

Subset Cryptanalysis

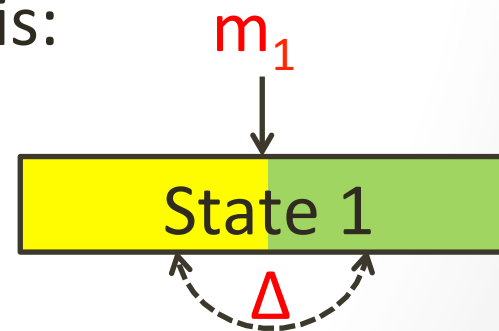
- In order to devise and analyze the attack we use a very common cryptanalysis framework which we call **subset cryptanalysis**
- Uses **subset characteristics** to track the evolution of subsets through the internal state of the cryptosystem
 - Associate a triplet (input subset, output subset, transition probability) to each internal operation

Internal Differential Cryptanalysis

- Introduced by Thomas Peyrin (Crypto 2010) in the analysis of Grostl
- Standard differential cryptanalysis:



- Internal differential cryptanalysis:



Generalized Internal Differential Cryptanalysis

- We generalize and extend it:
 - Shown to be applicable only to hash functions built using **separate data-paths**, whereas Keccak has only one **data-path**
 - The differences considered were between **2** parts of the state, whereas we consider more complex differential relations between **multiple parts of the state**

Internal Differences

Definitions

- In symmetric states all CSS's are equal
- In states which are almost symmetric the **differences** between the **first** CSS and the other **3** CSS's ($\Delta_1, \Delta_2, \Delta_3$) are of **low Hamming weight**
- We group all states with a **fixed** ($\Delta_1, \Delta_2, \Delta_3$) into an **internal difference set**

Internal Differences

Definitions

- Given a state u , the set $\{v \mid v=u+w \text{ and } w \text{ is symmetric}\}$ is an **internal difference set**
- The differences between the CSS's is specified by u which is a **representative state**
- A state v of a **lowest Hamming weight** defines the **weight** of the internal difference
- The **zero internal difference** contains the symmetric states and has a weight of 0

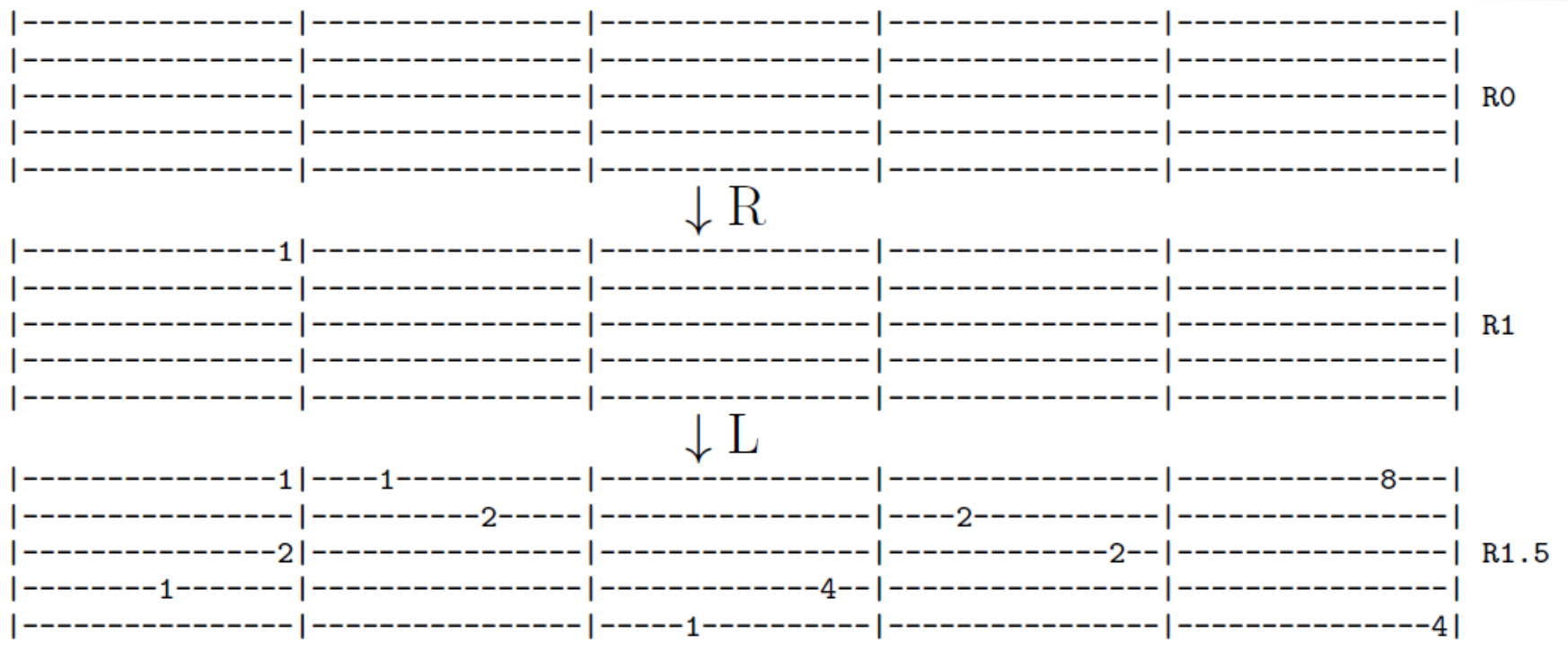
Internal Differential Characteristics

The Evolution of Internal Differences

- Any **symmetric** state chosen from the zero self-difference **remains symmetric** after applying Θ, ρ, π, χ
- Internal Differences are **affine subspaces**
- Their evolution through the **4** linear maps can be **easily tracked**

Internal Differential Characteristics

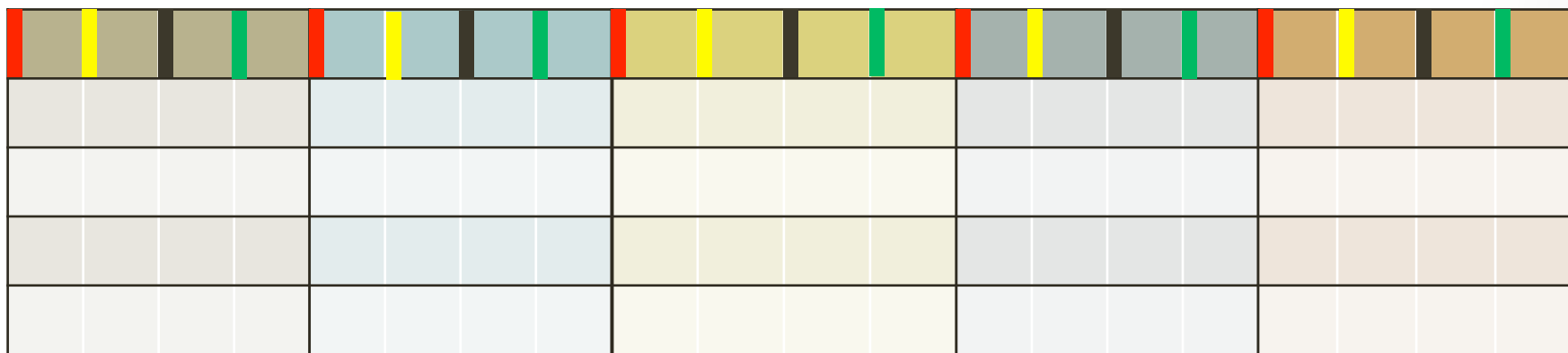
A 1.5-round Example



Internal Differential Characteristics

The Evolution of Internal Differences

- The evolution through χ is analyzed by considering **rotated row sets (RRS)**
- A RRS contains an Sbox (row) in the first CSS and its **3 symmetric counterparts** in the other CSS's



Internal Differential Characteristics

The Evolution of Internal Differences

- The input internal difference specifies the **differences** between the Sboxes of **each RRS**
- Each RRS can assume **exactly 32** values
- The distribution of the output internal difference can be **computed exhaustively**
- For **i=32** the output internal difference can be analyzed in a similar way to **standard differential cryptanalysis**

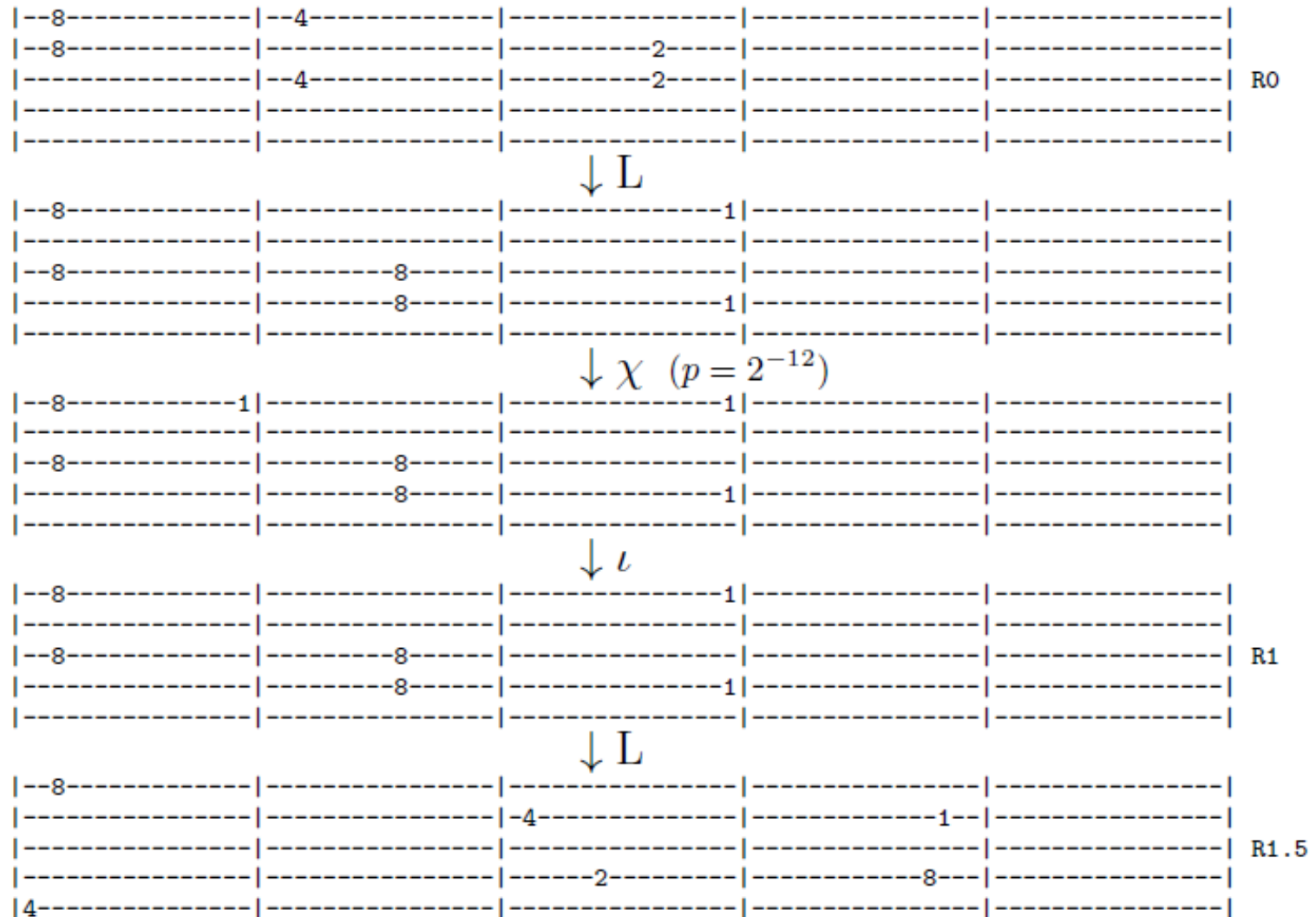
Internal Differential Characteristics

The Evolution of Internal Differences

- A RRS with a zero input difference is called **non-active**
- A non-active RRS passes through χ with probability **1**
- A **low-weight** internal difference passes through χ with high probability
- We look for internal differential characteristics composed of **low-weight** internal differences

Internal Differential Characteristics

Another 1.5-round Example



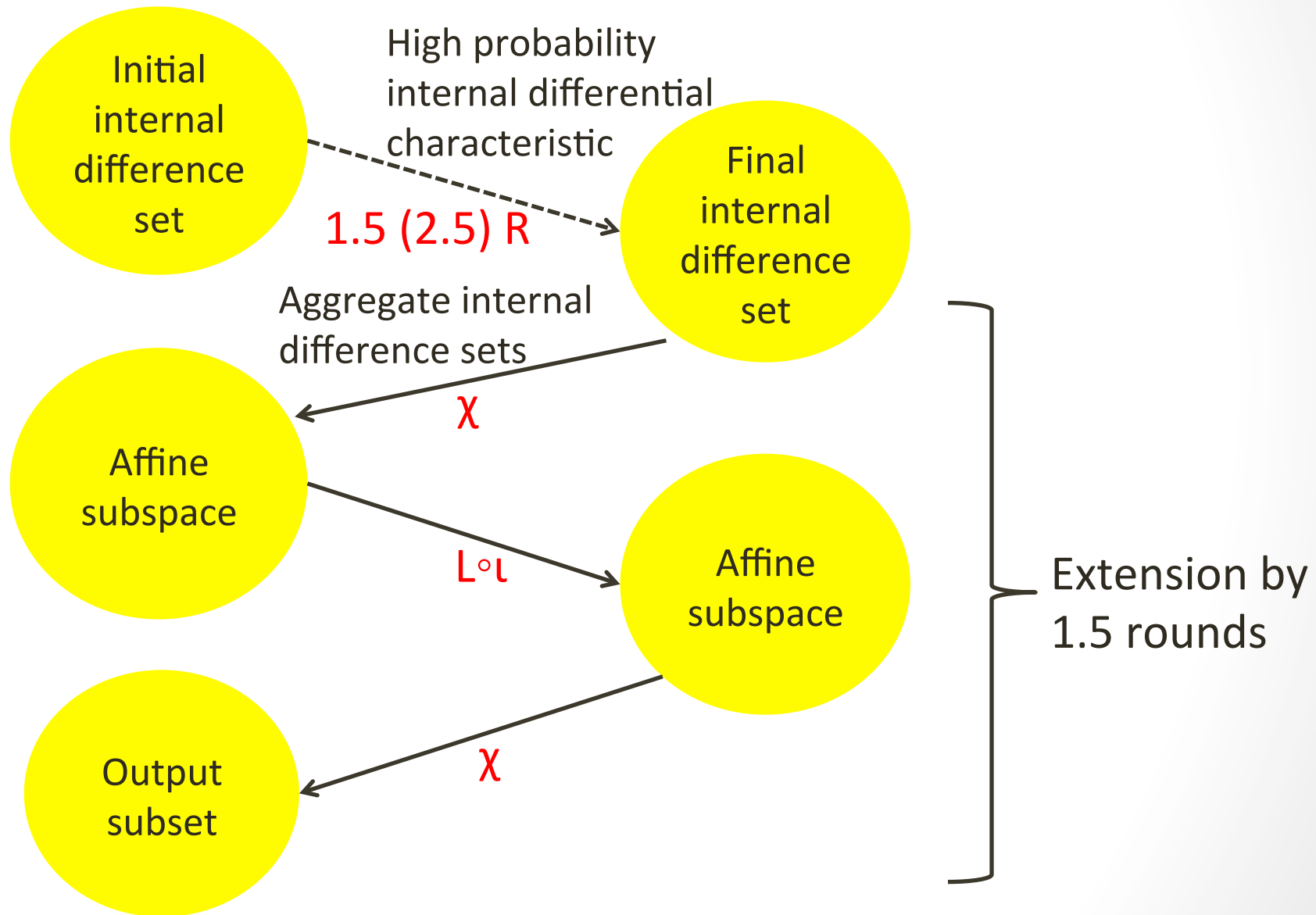
Extending Internal Differential Characteristics

- Given a characteristic that ends before the χ layer with a (relatively) **high weight** internal difference
- We extend it by **1.5 additional** rounds to a **subset characteristic**
 - **Do not restrict** its subsets to specific internal differences
 - **Avoid** the reduction in probability

Extending Internal Differential Characteristics

- Exploit the **low algebraic degree** of χ to **aggregate** all the possible internal differences at the output of χ to a single **affine subspace**
- Extend the characteristic by **1** round using an **affine subspace**
- Use the **low diffusion** of χ to bound the output subset size after an additional **0.5** round

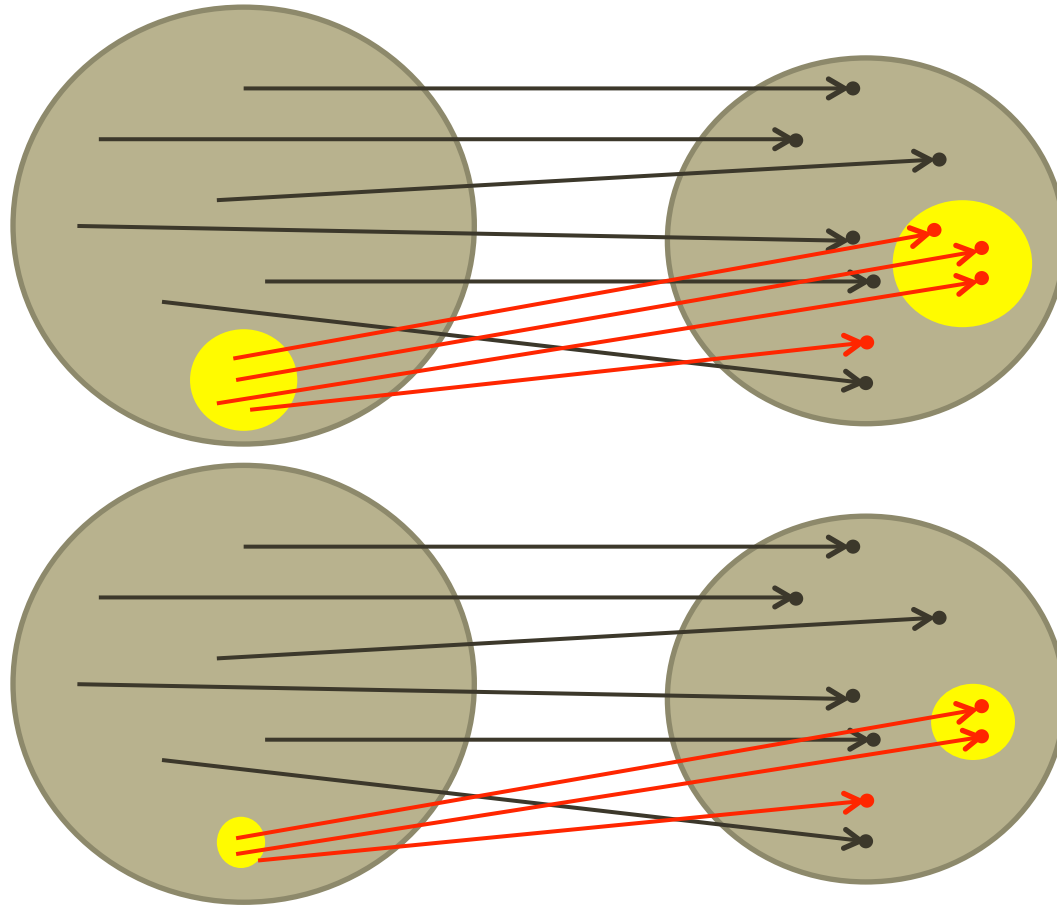
The Evolution of Subsets



Choosing the Rotation Index

- A **smaller** rotation index enforces **more symmetry relations**
 - **Reduces** the size of the **output** set
 - **Reduces** the size of the **input** set

Choosing the Rotation Index



- Choose the **smallest** value of $i \in \{1, 2, 4, 8, 16, 32\}$ for which the **input set** is **large** enough to find a **collision**

Collision Attacks

Practical Attacks

- A **3**-round collision in **Keccak-512** (with rotation index **i=4**)

M1=

88888888 88888888 66666666 66666666 AAAAAAAAA AAAAAAAAA 77777777
77777777 BBBB BBBB BBBB BBBB 11111111 11111111
88888888 88888888 CCCCCC CCCCCC

M2=

AAAAAAAA AAAAAAAAA 88888888 88888888 EEEEEEE EEEEEEE 99999999
99999999 99999999 99999999 99999999 99999999 88888888 88888888
CCCCCCC CCCCCC CCCCCC CCCCCC

Output=

56BCC94B C4445644 D7655451 5DD96555 71FA7332 3BA30B23 958408C5
64407664 41805414 11190901 6ABAA8BA A8ABAEFA 7EF8AEEE ECCE68DC
4EC8ACEC DD5D5CCC

Collision Attacks

Practical Attacks

- A 3-round collision in **Keccak-384** (with rotation index $i=4$)

M1=

FFFFFFFF FF7FFFFFFF BBBBBBBB BBFB BBBB 44444444 44444444 FFFFFFFFFF
FFFFFFFF 99999999 99999999 44444444 44C44444 44444444 44444444
44644444 44444444 AAAAAAAAA AAAAAAAAA 66666666 66666666 44444444
44444444 DDDDDDDD DD9DDDDD DDFDDDDD DDDDDDDD

M2=

33333333 33B33333 55555555 55155555 AAAAAAAAA AAAAAAAAA 77777777
77777777 44444444 44444444 66666666 66E66666 EEEEEEEE EEEEEEEE
11311111 11111111 CCCCCCCC CCCCCCCC FFFFFFFFF FFFFFFFFF 11111111
11111111 99999999 99D99999 DDFDDDDD DDDDDDDD

Output=

99999991 11199999 4440C444 405C60DC 00000000 0C100010 777677F7
73F77767 3550F597 55D57155 66666664 66666666

Collision Attacks

- The 2.5-round characteristic is used in a 4-round collision attack on **Keccak-384**
- The time complexity is 2^{147} (faster than the birthday bound by 2^{45})

Collision Attacks

- A **5**-round collision attack on **Keccak-256**
- Based on a **target internal difference algorithm**
 - An extension of the **target difference algorithm**
(FSE 2012)

The Target Internal Difference

Extending a Characteristic Backwards

BCE-1EDC-----	68A-49EE-----	4A6-2999-----	916-3-39-----	68A--D8-----
BCE-9E54-----	68A-49EA-----	4A6-2999-----	916-3-99-----	6CA--D8-----
FCE-9E54-----	68A-49EE-----	4A2-2999-----	916-3-99-----	68A--D8-----
BCE-9E54-----	68A-49EE-----	4A6-2999-----	916-3-99-----	68A--D8-----
BCE-9E54-----	68A-49EE-----	4A6-2B99-----	916-3-99-----	69A--D8-----

↑ R ($p=2^{-16}$)

-----8-----	---4-----	-----2-----	-----	-----4-----
-----8-----	---4-----	-----2-----	-----	-----4-----
-----	-----	-----	-----	-----
-----	-----	-----	-----	-----
-----	-----	-----	-----	-----

↓ L

-----8-----	---4-----	-----	-----	-----
-----	---4-----	-----	-----	-----
---8-----	-----8-----	-----	-----	-----
---2-----	-----8-----	-----	-----	-----
8-----	-----	-----	-----	-----

↓ χ ($p = 2^{-16}$)

-----8-----	---4-----	-----	-----	-----
-----	---4-----	-----	-----	-----
---8-----	-----8-----	-----	-----	-----
---82-----	-----8-----	-----	-----	-----
8-----	-----	-----	-----	-----

↓ ι

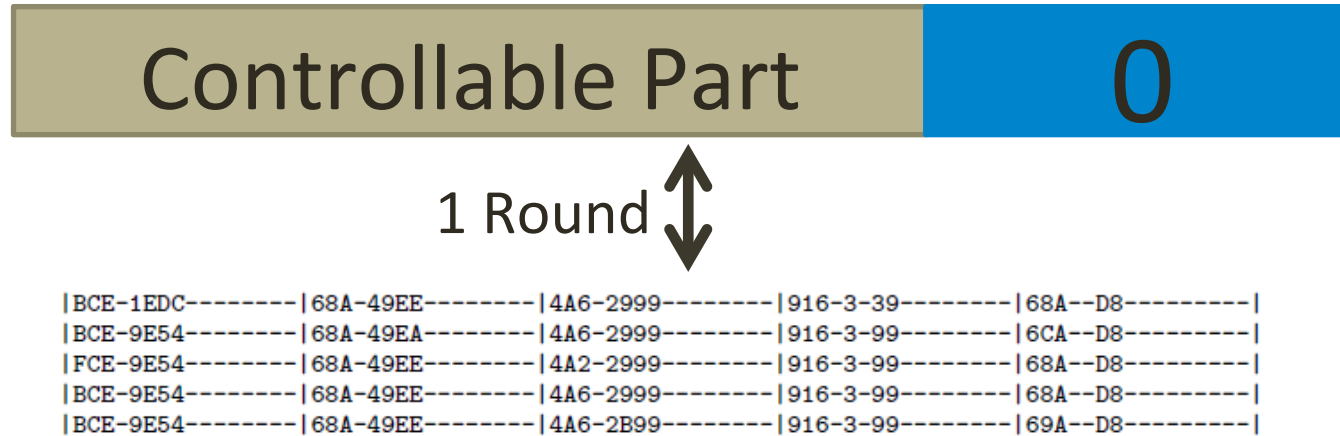
8--8-82-----	---4-----	-----	-----	-----
-----	---4-----	-----	-----	-----
---8-----	-----8-----	-----	-----	-----
---82-----	-----8-----	-----	-----	-----
8-----	-----	-----	-----	-----

↓ L

8--8-82-----	-----4-----	-----	-----	-----
-----	-----	---4-----	---1-----	-----
---8-----	-----	-----	-----	---2-----
-----	-----	---2-----	-----	-----
-----	-----	-----	---1-4-----	-----

The Target Internal Difference Algorithm

Linking a Characteristic Form an Initial State



- The time complexity of the optimized attack is about 2^{115}
 - Faster than the birthday bound by 2^{13}

Conclusions and Future Work

- We presented the **first collision attacks** on round reduced **Keccak-384** and **Keccak-512**
 - Some of them are practical
- For **Keccak-256** we **increased** the number of rounds that can be attacked from **4** to **5**
 - We are still **very far** from attacking the full **24** rounds
- An interesting future work item is to find **better** internal differential characteristics for Keccak **or** to **prove** that they do not exist

Thank you for your attention!